

# Stochastic Differential Games and associated Bellman-Isaacs Equations

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## An Introduction

### Deterministic differential games:

Rufus Isaacs (1914 - 1981) :

the first to study differential games:  
1951 (published: 1958):

a pursuit game, more precisely  
"the homicidal chauffeur problem":

two players:

P1: a runner (rather slow, but highly manoeuvrable)  
against

P2: a driver of a motor vehicle (much faster,  
but less manoeuvrable)

Objective of P2: to flatten down the pedestrian,  
Objective of P1: to elude the car

for the solution was used:

level set method, variational calculus.

Rufus Isaacs considered special cases.

Since the pioneering work by Rufus Isaacs:  
 a lot of works on deterministic and also  
 stochastic differential games:

Friedman (1971)

Elliot / Kalton (1972)

Fleming

Fleming / Souganidis (1989)

Marc Quincampoix (GDR "Games")

Pierre Cardaliaguet / Catherine Rainer

An easy model of a pursuit - evasion game:

Evasion (P1):  $\dot{x}(t) = u_t, t \geq 0, u_t \in U, x_0 \in \mathbb{R}^2$

Pursuer (P2):  $\dot{y}(t) = v_t, t \geq 0, v_t \in V, y_0 \in \mathbb{R}^2$

+  $\rho \geq 0$ ;

Objective of P1:  $\inf_{t \geq 0} |x(t) - y(t)| \geq \rho$

Objective of P2:  $\inf_{t \geq 0} |x(t) - y(t)| \leq \rho$ .

+ If it is not for catching the evader but a competition for money:

Example:

$J(u, v) = |x(T) - y(T)|$  amount P2 has to  
 pay to P1.

The dynamics can be written as:

$$d\begin{pmatrix} x(t) \\ y(t) \end{pmatrix} = \begin{pmatrix} u_t \\ v_t \end{pmatrix} dt, t \geq 0; \begin{pmatrix} x_0 \\ y_0 \end{pmatrix} \in \mathbb{R}^4.$$

More general, for initial data  $(t, x) \in [0, T] \times \mathbb{R}^d$ :

$$\begin{cases} dX_s^{t,x,u,v} = b(X_s^{t,x,u,v}, u_s, v_s) ds + \sigma(X_s^{t,x,u,v}, u_s, v_s) dw_s, \\ X_T^{t,x,u,v} = x \end{cases}$$

payoff / cost functional:

$$J(t, x; u, v) = E \left[ \Phi(X_T^{t,x,u,v}) + \int_t^T f'(X_s^{t,x,u,v}, u_s, v_s) ds \right]$$

2-persons zero-sum game:

$$\begin{aligned} P_1: J(t, x; u, v) &\rightarrow \max \text{ over } u \in U_{t,T} \\ P_2: J(t, x; u, v) &\rightarrow \min \text{ over } v \in V_{t,T} \end{aligned}$$

Which type of game? Control against control?

In special cases possible:

+ linear-quadratic SDE:

existence of a saddle point control  
 $(u^*, v^*) \in U_{t,T} \times V_{t,T}$ ,

characterized by the associated Riccati equation.

+ S. Peng / J.-P. Lepeltier / S. Hamadène (1997  
 and consecutive works):

If

- $\tilde{G}(x, u, v) = \tilde{G}(x)$ ,  $\exists \tilde{G}(x)^{-1} \in R^{d \times d}$ ,  $|\tilde{G}(x)|^{-1} \leq c$ ,
- Cost functional given by a BNE with controls,
- Isaacs condition holds

then J. Peng's Comparison theorem for BNEs shows:  $\exists$  saddle point control  $(u^*, v^*)$ .

$(u^*, v^*) \in U_{t,T} \times V_{t,T}$  is saddle point  
iff

$$\mathcal{J}(t, x, u, v^*) \leq \mathcal{J}(t, x, u^*, v^*) \leq \mathcal{J}(t, x, u^*, v)$$

Note: If there is a saddle point:

$$\boxed{\begin{aligned} W(t, x) &:= \sup_{u \in U_{t,T}} \inf_{v \in V_{t,T}} \mathcal{J}(t, x, u, v) \\ &= \inf_{v \in V_{t,T}} \sup_{u \in U_{t,T}} \mathcal{J}(t, x, u, v) \end{aligned}}$$

$\hat{=}$  game has a "value"

(value doesn't depend on who begins the game).

In general: no saddle point, even no value,  
although Isaacs condition holds.

Example:

A (simplified) pursuit game:

P1 (rabbit):  $\dot{x}_1^u(s) = u_5, s \in [0, T], u_5 \in \overline{B_1(0)} \subset R^2$   
P2 (wolf):  $\dot{x}_2^v(s) = v_5, s \in [0, T], v_5 \in \overline{B_2(0)} \subset R^2$   
 $x_1^u(0) = x^1, x_2^v(0) = x^2 \in R^2$

Rabbit:  $\mathcal{U} = L_A^0(0, T; \overline{B_1(0)})$ ,  
Wolf:  $\mathcal{V} = L_A^0(0, T; \overline{B_2(0)})$ .

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"lunch time"  $T$ :  $x = (x^1, x^2)$ :

Rabbit:  $J(0, x; u, v) = E[T | X_1^u(T) - X_2^v(T)|]$   
 Wolf:  $J(0, x; u, v) \rightarrow \max_{u \in \mathcal{U}}$  over  $v \in \mathcal{V}$   
 $\min_{v \in \mathcal{V}}$  over  $u \in \mathcal{U}$

$w(0, x) := \sup_{u \in \mathcal{U}} \inf_{v \in \mathcal{V}} J(0, x; u, v)$   
 $u(0, x) := \inf_{v \in \mathcal{V}} \sup_{u \in \mathcal{U}} J(0, x; u, v)$ .

1) Rabbit begins the match:  
Suppose  $|x^1 - x^2| \leq T$ .

Rabbit chooses:  $u \in \mathcal{U}$ ,

Wolf can then choose:  $v_5 = \frac{x^1 - x^2}{T} + u_5, s \in [0, T]$ .  
 $\Rightarrow v \in \mathcal{V}$  and

$J(0, x; u, v) = 0$ , ie.,  $w(0, x) = 0$ .

2) Wolf begins the match:

wolf chooses  $v \in V$ ,

rabbit then can choose

$$u_5 := \begin{cases} 0, & \text{if } E[x^1 - x_2^v(T)] \geq \frac{(T_{11})^2}{2}, \\ (T_{11})e, & \text{if } E[x^1 - x_2^v(T)] < \frac{(T_{11})^2}{2}, \end{cases} \quad a)$$

$s \in [0, T]$ , for  $e \in \mathbb{R}^2$ ,  $\|e\| = 1$ .

{One could take  $e = \frac{E[x^1 - x_2^v(T)]}{\|E[x^1 - x_2^v(T)]\|}$ ,  
but any other direction also works}

Then:  $u \in C^0(0, T; \overline{B_1(0)})$

$\subset C_F^0(0, T; \overline{B_1(0)}) = U$  admissible control

Moreover:

$$\begin{aligned} J(0, x, u, v) &= E[(X_1^u(T) - X_2^v(T))] \\ &= \begin{cases} E[x^1 - x_2^v(T)] \geq \frac{(T_{11})^2}{2}, & \text{in case a)} \\ E[x^1 + (T_{11})T/e - x_2^v(T)] \\ \geq (T_{11})T/e - \frac{(T_{11})^2}{2} \geq \frac{(T_{11})^2}{2}. \end{cases} \end{aligned}$$

$$\Rightarrow \sup_{u \in U} J(0, x, u, v) \geq \frac{(T_{11})^2}{2}, \forall v \in V.$$

$$U(0,x) = \inf_{v \in V} \sup_{u \in U} J(0,x,u,v) \geq \frac{(T_{11})^2}{2}. \quad I7$$

$$\Rightarrow \boxed{W(0,x) = 0 < \frac{(T_{11})^2}{2} \leq U(0,x),} \\ \forall x = (x^1, x^2) \text{ with } |x^1 - x^2| \leq T.$$

Concept used by  
Fleming / Souganidis (1989) in their pioneering  
paper on SDEs:

Strategies - controls.

In our mini-course: generalisation of  
of their paper

Remark:

P. Cardaliaguet / C. Rainer:

- Nonanticipating strategy with delay
- Contre - nonanticipating strategy, with delay.